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Math 403 Fall 2011 UMD

Lecture 1

Functions, Sets, Notations

§0 Sets of numbers and notation for them

- $\mathbb{N} = \{0, 1, 2, 3, \dots\}$ = set of natural numbers
- $\mathbb{Z} = \{\dots, -3, -2, -1, 0, 1, 2, \dots\}$ = set of integers
- $\mathbb{Q} = \left\{ \frac{a}{b} : a, b \in \mathbb{Z}, b \neq 0 \right\}$ = {rational nos}
- \mathbb{R} = {set of real nos}

~~not~~

If S is a set,

We write $a \in S$ to mean that a is an element of S ,

$a \notin S$ means that a is not an element of S

Ex $-1 \in \mathbb{Z}$, $-1 \notin \mathbb{N}$, $\pi \in \mathbb{R}$, $\pi \notin \mathbb{Z}$

We write $A \subseteq B$ to mean that every element of A is in B . If $A \subseteq B$ then A is said to be a sub-set of B .

Ex $\mathbb{N} \subseteq \mathbb{Z} \subseteq \mathbb{Q} \subseteq \mathbb{R}$

§1 Defining sets by properties

If S is a set and P is a property that an element of a set may or may not have

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then we can use P to define a subset T consisting of all elements of S satisfying P .
We use set-builder notation to write this set

$$T = \{x \in S : P(x)\}$$

Ex $E := \{n \in \mathbb{Z} : n \text{ is even}\} = \{\text{even integers}\}$

$$O := \{n \in \mathbb{Z} : n \text{ is odd}\} = \{\text{odd integers}\}$$

$$[0, \infty) := \{x \in \mathbb{R} : x \geq 0\}$$

$$\mathbb{Z}_+ := \{n \in \mathbb{Z} : n > 0\}$$

Ex 2 logical symbols and notation If P, Q properties, then

$P \Rightarrow Q$ means P implies Q

$P \Leftarrow Q$ means Q implies P

$P \Leftrightarrow Q$ means P iff Q .

[$P \vee Q$	means either P or Q hold]	leave off
	$P \wedge Q$	means both P and Q hold		
	$\neg P$	means P doesn't hold.		

\exists means there exists

~~$\forall x \in [0, \infty) \Rightarrow x = y^2$ for some~~

\forall means \forall

Ex $x \in [0, \infty) \Rightarrow \exists y \in \mathbb{R}, x = y^2$

That's a shorthand for saying any non-negative real no is a square.

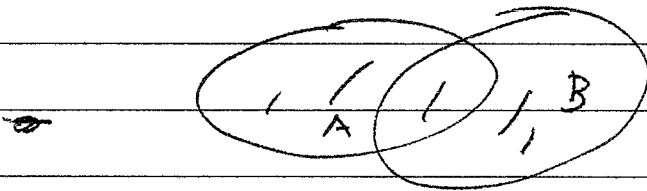
§3 Operations on Sets

$A \cup B = \{x : x \in A \text{ or } x \in B\}$, Union

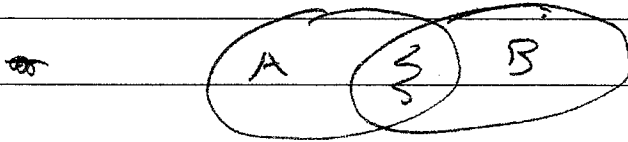
$A \cap B = \{x : x \in A \text{ and } x \in B\}$, intersection

$A \setminus B = \{x \in A : x \notin B\}$, complement difference

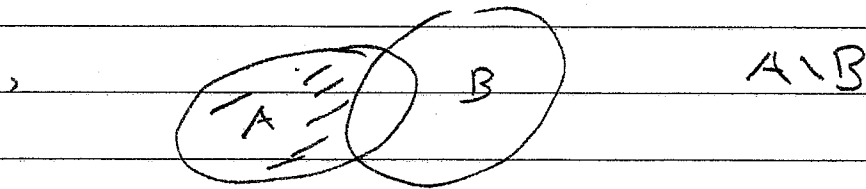
If $B \subseteq A$ then $A \setminus B$ is called the complement of B in A .



$A \cup B$



$A \cap B$



$A \setminus B$

~~Ex 2~~ $A \setminus B =$

There is a special set called the empty set which has no elements. We write it as \emptyset .

We say A and B are disjoint if $A \cap B = \emptyset$.

Ex $E \cup O = \mathbb{Z}$ $E \cap O = \emptyset$
Complement of E in \mathbb{Z} is O

Q4 Other operations

Power Set If S is a set

$$P(S) = \{ \text{all subsets of } S \}$$

Ex $S = \{1\}$, $P(S) = \{ \emptyset, \{1\} \}$
 $S = \{1, 2\}$, $P(S) = \{ \emptyset, \{1\}, \{2\}, \{1, 2\} \}$
 $S = \{1, 2, 3\}$

Prmk If S is finite with n elts, then $P(S)$ has 2^n elts

Cartesian ~~star~~ product If S, T are sets then $S \times T$ is set of all ordered pairs (s, t) where $s \in S, t \in T$.

Ex $S = \{1, 2\}$, $T = \{4, 5, 6\}$

$$S \times T = \{ (1, 4), (1, 5), (1, 6), (2, 4), (2, 5), (2, 6) \}$$

If S, T finite with n, m elts resp. then $S \times T$ has nm elts.

Remarks ① $(s, t) = (s', t') \Leftrightarrow s = s', t = t'$

② In math officially everything is a set. (But we don't really like to think that way.)

If s, t are sets then by def

$$(s, t) := \{\{s\}, \{s, t\}\}$$

It is ~~is~~ a set with two elts:

- $\{s\}$

- the set $\{s, t\}$ having two elts: s and t .

This is confusing enough that we try not to bring it up but you can check that, with this def,

$$(s, t) = (s', t') \Leftrightarrow s = s', t = t'$$

③ Even the elts of \mathbb{N} are formally sets. So the way it's officially set up

$$0 = \emptyset \quad \text{and if } n \in \mathbb{N} \text{ is defined}$$

$$n+1 = n \cup \{n\}$$

$$\text{So } 1 = \{\emptyset\}, \quad 2 = \{\emptyset, \{\emptyset\}\}$$

$$3 = \{\emptyset, \{\emptyset\}, \{\emptyset, \{\emptyset\}\}\}$$

This is super confusing so we never think of it this way. Don't write \emptyset for 0 when you mean 0

as a natural number!

But notice

0 has	0 elts
1 has	1 elt
2 has	2 elts

~~Mappings~~ Suppose S, T are sets. They ~~are~~ informally speaking a function from S to T is

§5 Relations and mappings

Suppose S, T are sets. Then a relation between S and T is a subset $R \subseteq S \times T$. If $s \in S, t \in T$ then t is related to s if $(s, t) \in R$.

A mapping from S to T is a relation $f \subseteq S \times T$ st for each $s \in S$ there exists a unique $t \in T$ st $(s, t) \in f$. If $(s, t) \in f$ then, we write $f(s)$ ~~$f(s)$~~
 $f(s) = t$. This defines a rule associating to each element of s a unique element of t .

Mappings are also called functions. Often we define the function by giving the rule.

Ex $S = \mathbb{R}, T = \mathbb{R}$

$$f(x) = x^2$$

This is a rule taking any real number to its

square. Officially $f = \{ (x, y) \in \mathbb{R} \times \mathbb{R} : y = x^2 \}$.
 But we like thinking of f better as a rule, more
 than as a subset of $\mathbb{R} \times \mathbb{R}$. ~~Sometimes we call~~

We write $f: S \rightarrow T$ to indicate that f is a mapping
 from S to T . Sometimes we give the rule by writing
 $s \mapsto f(s)$ or $f(s) = \text{"some rule"}$.

Ex. $f: \mathbb{Z} \rightarrow \mathbb{Z}$ given by
 $n \mapsto 2n$

This is just another ~~and~~ way of saying $f(n) = 2n$.

• $f: \mathbb{N} \rightarrow \mathbb{N}$ given by $n \mapsto n!$

Obs If $f: S \rightarrow T$, $f': S \rightarrow T$ are two functions
 and $f(s) = f'(s)$ for all $s \in S$. Then
 $f = f'$. / Intrinsically, f 's def by rule

A special function If S is a set then $\text{id}_S: S \rightarrow S$,
 etc. stupid f 's given by $s \mapsto s$.

Composition of Functions Suppose $f: S \rightarrow T$, $g: T \rightarrow U$
 are f 's. We can define a new function
 $h = g \circ f$ called the composition of g with f
 or g composed with f by the rule

$$h(s) = g(f(s)).$$

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Thm Suppose $f: S \rightarrow T$, $g: T \rightarrow U$, $h: U \rightarrow V$ are f'ns. Then

$$(i) \quad g \circ \text{id}_T = \text{id}_U \circ g = g$$

$$(ii) \quad h \circ (g \circ f) = (h \circ g) \circ f$$

Pf (a) $(g \circ \text{id}_T)(t) = g(\text{id}_T(t)) = g(t)$

So $g \circ \text{id}_T = g$. Similarly $\text{id}_U \circ g = g$.

$$(b) \quad (h \circ (g \circ f))(t) = h((g \circ f)(t))$$

$$= h(g(f(t))) = (h \circ g)(f(t))$$

$$= ((h \circ g) \circ f)(t).$$

Since this holds for all $t \in T$, $h \circ (g \circ f) = (h \circ g) \circ f$.

Range, onto, into

Suppose $f: S \rightarrow T$ is a function and $A \subseteq S$. We write

$$f(A) = \{ f(a) : a \in A \}$$

This is called the image of A , $f(S)$ is ~~the~~ called the range of f , f is onto if $f(S) = T$.

f is said to be 1-1 if $f(s) = f(s') \Rightarrow s = s'$.

Ex - The function $f: \mathbb{R} \rightarrow \mathbb{R}$ given by $x \mapsto x^2$ is not onto, not 1-1. $f(\mathbb{R}) = [0, \infty)$.
Not 1-1 b/c $f(1) = f(-1)$.

- The function $f: \mathbb{R} \rightarrow [0, \infty)$ given by same rule $x \mapsto x^2$ is onto but not 1-1

- The fn $f: [0, \infty) \rightarrow \mathbb{R}$ given by $x \mapsto x^2$ is not onto but is 1-1.

Def A function $f: S \rightarrow T$ is said to be a 1-1 correspondence if it is both 1-1 and onto. In this case we can define a function $g: T \rightarrow S$ given by

$$g(t) = \text{the unique } s \in S \text{ st } f(s) = t.$$

We then have $g \circ f = \text{id}_S$

$$f \circ g = \text{id}_T.$$

g is sometimes called the inverse of f .

Ex $f: (0, \infty) \rightarrow (0, \infty)$ given by
 $x \mapsto x^2$

is 1-1, onto.

The inverse is $g: f^{-1}(y) = \sqrt{y}$.